Last time

Classical cryptography ⇒ one-time pad

Block ciphers

Block cipher modes and MACs

Today: cryptographic hash functions

MAC Requirements

- 1. Arbitrary-length message
- 2. Small, fixed MAC length
- 3. Computationally efficent
- 4. Collision resistance:
 - o can't generate another message with the same MAC
 - o can't generate another message with any valid MAC

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Note: the Sealed Authenticator System (SAS) codes on a nuclear-armed submarine probably don't use keyed MACs, but rather purely-random codes that no human eyes have ever seen. Source: Waller, "Practicing for Doomsday", Time Magazine, 4 Mar 2001.

MAC generalization

What if we don't need a block cipher?

What if we don't want to use a key?

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But why wouldn't we want to use a key?

AAA[A]

Category	Question
Authentication	Is something/someone authentic (is it really you)?
Authorization	Are you allowed to do that?
Accounting	Who has used which resources?
Audit	Who did what to what?

Message authentication vs principal authentication

Examples of	include the auth	include the authenticated orders in Crimson	
Tide and the payment authorization messages describde by the EMV protocol. In both of these			
cases, there are	required besides the	itself.	
When authenticating	instead of	, we can use messages	
in which the message itself is the secret, for example			

Passwords

Old and terrible, but...

Dictionary attack

- online
- offline ???

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We'll talk later in the term about protocols that we can use for authentication based on a third party, but at some point, *somebody* has to store a password

A dictionary attack is a brute-force attack: instead of trying every possible key for a cipher, you try every possible password from a dictionary. This is generally cleverer than trying "aaaaaa", "aaaaab", etc., as some passwords are (unfortunately) likelier to be chosen than others. Also, the dictionary may include more than just "dictionary" words!

Threats to authentication

External threats

- password guessing
- MAC-based challenge/response guessing human-computable?

Internal threats

- password database could be stolen
- ... but so could a secret key for validating MACs!

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MAC-based schemes only work when the secret key ______. We can't guarantee that in general-purpose computers.

We'll talk later about public-key schemes that can help with the theft issue, but they don't help with the human-computability problem.

Cryptographic hash functions

Remember hash tables' hash functions?

- variable-length input
- fixed-length output

Cryptographic hash functions

MD4, MD5, SHA-1, RIPEMD-160, Whirlpool, SHA-2 (SHA-224, SHA-256, SHA-384, SHA-512, SHA-512/224, SHA-512/256), SHA-3, BLAKE2/3...

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These properties sound like some of the properties of MACs: variable-length input, fixed-length output, computationally efficient and avoiding collisions. However, while regular hash functions try to avoid collisions, they do happen, because the consequences of a collision aren't terribly serious. If we start to see lots of collisions in a hash table, we can always increase the size of the table.

Cryptographic hash functions, however, are something entirely different. A cryptographic hash function should still be fairly efficient to compute (in practice, we can hash millions of MB/s), but efficiency has to be traded off for *much* stronger ________. Once we start sending messages around with cryptographic hashes, we can't recall all of the messages and re-hash them. Instead, we must be very strict about up front.

Cryptographic hash function

Diffusion: small changes ⇒ large effects

All values should be equally likely

Should resist:

Collision attack: find X_1,X_2 s.t. $h(X_1)=h(X_2)$ Preimage attack: given $h(X_1)$, find X_2 s.t. $h(X_1)=h(X_2)$ 2nd preimage attack: given X_1 , find $X_2 \neq X_1$ s.t. $h(X_1)=h(X_2)$

Collision attack			
Finding	that hash to the same value. When we get to digital		
	ion attacks can be quite important: if you can generate two		
messages with different meanings but the same hash, you can cause a lot of trouble! However,			
such attacks aren't so useful for	r password security.		
Even with the strongest hash fu	unction, collisions are		
due to the birthday	paradox. However, "easier" doesn't have to be "easy": if the hash		
output is large, you can still ha	ve a lot of work to do! $\sqrt{2^n}$ can still be a large number if n is		
big enough			
Preimage attack			
Finding an input that hashes to the same value as a given hash. This could be the same input			
that was originally used to gene	erate the hash or a different one.		
Second preimage attack			
Finding ai	nput that will hash to the same value as a given input. This is		
like a collision attack, but much harder: instead of generating lots of messages and finding two			
that hash to the same value, you have to find one that hashes to the same value			

Password hashing

What does this have to do with passwords?

Resisting offline dictionary attacks*

Rainbows† and salt

Iterative password hashing (KDFs)

† Oeschslin, "Making a Faster Cryptanalytic Time-Memory Trade-Off", CRYPTO 2003: Advances in Cryptology - CRYPTO 2003, 2003. DOI: 10.1007/978-3-540-45146-4_36.

We don't need any cryptography to resist an online dictionary attack. Protecting password				
databases is, instead, all about resisting	, where an adversary has			
gained access to a password database and they want to get passwords from it. Without any				
cryptography, they can simply do a database lookup. With cryptography, however, we can				
make things much harder for them.				
As a (very bad!) alternative to password hashing, check out this analysis of a major password				
database breach at Adobe.				
Tools like GPUs are really good at parallel computation. Attackers can use them to try lots and				
lots of passwords concurrently to see if they can find the correct one (a bit like the Bombes in				
Bletchley Park!).	(KDFs) make life harder for an attacker			
by forcing computation to be	. There is a cost for the user, too, but it's			
insignificant compared to the benefit of not having your password cracked when a business				
suffers a data breach!				

^{*} see, e.g., John the Ripper

What makes a good password?

(we'll answer this next time)

MAC generalization

What if we don't want to use a key?

What if we don't use a block cipher?

HMAC: hash-based message authentication code*

$$h\left((k \oplus p_o)||h((k \oplus p_i)||text)\right)$$

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An HMAC uses a hash function *with* a key. This provides the same security properties as a block-cipher–based MAC, just with a different underlying cryptographic algorithm. HMACs are pretty popular in circumstances where you'd be doing a bunch of hashing anyway (e.g., Transport Layer Security cipher suites, which we'll talk about later).

^{*} Bellare, Canetti and Krawczyk, "Keying Hash Functions for Message Authentication", CRYPTO 1996, 1996. Standardized by NIST (FIPS 198-1) and the IETF (RFC 2104).