## **Recall**

### AAA[A]

### **Passwords**

- "password" files
- dictionary attacks: online vs offline

## **Cryptographic hash functions**

### Password hashing

Resisting offline dictionary attacks\*

Rainbows† and salt

Iterative password hashing (KDFs)

† Oeschslin, "Making a Faster Cryptanalytic Time-Memory Trade-Off", CRYPTO 2003: Advances in Cryptology - CRYPTO 2003, 2003. DOI: 10.1007/978-3-540-45146-4 36.

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We don't need *any* cryptography to resist an online dictionary attack. Protecting password databases is, instead, all about resisting \_\_\_\_\_\_\_\_, where an adversary has gained access to a password database and they want to get passwords from it. Without any cryptography, they can simply do a database lookup. With cryptography, however, we can make things much harder for them.

As a (very bad!) alternative to password hashing, check out this analysis of a major password database breach at Adobe.

Tools like GPUs are really good at parallel computation. Attackers can use them to try lots and lots of passwords concurrently to see if they can find the correct one (a bit like the Bombes in Bletchley Park!). \_\_\_\_\_\_\_ (KDFs) make life harder for an attacker by forcing computation to be \_\_\_\_\_\_\_. There is a cost for the user, too, but it's insignificant compared to the benefit of not having your password cracked when a business suffers a data breach!

<sup>\*</sup> see, e.g., John the Ripper

### What makes a good password?

#### Hard to guess

#### Complex?



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Fundamentally, it should be hard for an attacker to guess a password.

We have ideas about what makes guessing harder: not using common words, maybe making passwords long, maybe using funny symbols. Some of these ideas are intuitive, others have been \_\_\_\_\_\_\_\_. But why would those things make it harder to guess a password?

Before talking about sensible password policies, we need to understand just a little bit of information theory. One way of describing hard-to-guess-ness — but one which is often misunderstood — is the information-theoretic concept of *entropy*.

### **Entropy**

#### A measure of information

• or disorder, or chaos...

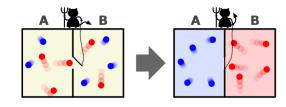
• thermodynamics: Maxwell's demon

• units: Shannons (a.k.a., bits!)



The concept of entropy is much older than computing: entropy has been used for a long time in thermodynamics to describe the disordered-ness of systems. In a closed system, entropy is a monotonically non-decreasing quantity, i.e., left alone, a closed system will become less ordered and more chaotic. The "heat death of the universe" doesn't refer to things getting really hot, it refers to an increase of random motion to the point that there is no structure, no separation between hot and cold, so no useful work can be done.

Maxwell's demon refers to a theoretical idea that links thermodynamics and information theory. Normally, if you bring hot and cold things together, their temperature evens out. However, if you could control the interface between two



chambers of gas such that you open the door for faster molecules going right and slower molecules going left, but close the door for other molecules, you could make lukewarm gas turn into hot gas in one chamber and cold in the other. Would this violate the second law of thermodynamics? No, because the \_\_\_\_\_\_ affects the entropy!

### Measuring entropy

#### Shannon entropy:

$$H(\mathbf{X}) = -\sum_{i=0}^n P(x_i) \log_b P(x_i)$$

### Hartley function:

$$H_0(\mathbf{X}) = \log_b |\mathbf{X}|$$

Guessing entropy, which is more directly relevant to password guessing, is a bit harder to calculate, but it is bounded by Shannon entropy. See: Massey, "Guessing and Entropy", Proc. IEEE Int. Symp. on Info. Th., 1994.

Claude Shannon defined one way of measuring entropy, based on how \_\_\_\_\_\_ specific values in a distribution are. In this definition, a distribution that includes only equally-likely values will have higher entropy than a distribution that has a different distribution.

Important note: \_\_\_\_\_\_

The Hartley function can be used to compute Shannon entropy \_\_\_\_\_\_\_

## Estimating password entropy

Eight random alphanumeric characters:

$$H_0(\mathbf{X}) = \log_b |\mathbf{X}| = \log_2 |36^8| = 41.6 \; \mathrm{Sh} \; \mathrm{(bits)}$$

Four diceware\* words + two numbers:

$$H_0(\mathbf{X}) = \log_2 \left| \left( 6^4 
ight)^4 imes 10^2 
ight| = 4 imes \log_2 \left| 6^4 
ight| + \log_2 \left| 10^2 
ight| = 48.0 ext{ Sh}$$

But people don't choose random passwords!

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In this example, we're using 26 possible letters plus 10 digits, so a total of 36 symbols that can be used. Eight of these symbols means that there are  $36^8$  possible passwords that can be formed with this scheme; if they are all \_\_\_\_\_\_, we can use the Hartley function to calculate the entropy of this \_\_\_\_\_\_.

These calculations show how we can work with distributions that contain multiple components, e.g., multiple words or multiple classes of characters. It's worth noting, however, that the entropy of the distribution of passwords derived from just four diceware words would be:

$$H_0(\mathbf{X}) = \log_2 \left| \left( 6^4 
ight)^4 
ight| = 4 imes \log_2 \left| 6^4 
ight| = 41.4 ext{ Sh}$$

So, four diceware words will stand up to a brute-force attack about as well as a random eight-character alphanumeric password. But which will be easier for a human to remember?

Calculating the entropy of random distributions isn't so hard. Unfortunately, however, people typically don't use randomly-generated passwords (although they should!).

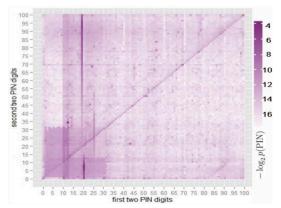
<sup>\*</sup> See EFF instructions plus Bonneau, "Deep Dive: EFF's New Wordlists for Random Passphrases", EFF, 2016.

# Actual password entropy

People don't choose random passwords!

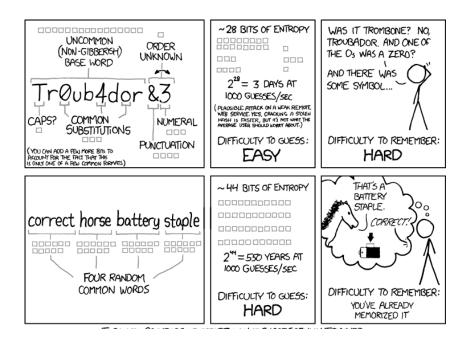
#### Or even random PINs!

- diagonal line: repetition
- lots of 19xx and even 20xx



Wang et al., "Understanding Human-Chosen PINs: Characteristics, Distribution and Security", in *ASIA CCS '17*, 2017. DOI: 10.1145/3052973.3053031.

This graphic was taken from PINs that were found as a subset of the password in the original RockYou data set. For those curious about "RockYou 2021": here's a nice writeup.



XKCD, although it's a web comic, often has spot-on analysis. Comedy and satire can make a point in a punchier way than a prosey explanation, and this is no exception!

### Entropy vs password strength

- Shannon, Hartley entropies have clear definitions
- Other entropies: min-entropy, guessing entropy...
- Entropy a measure of a distribution, not a single password
- Can estimate password entropy assuming random selection
- One website with one breach... haveibeenpwned.com

### What can we learn from this?

If we assume that	, we can compute the	
entropy of	. You should note that's	
exactly how I've phrased one of the questions in Assignment 2!		
We can learn things about actual user password choices, etc., from password databases that		
have been leaked. We can also learn whether or not particular passwords have been		
compromised!		
If a password has been leaked in the clear,		
Perhaps they weren't hashing, perhaps they weren't salting, perhaps they were		
allowing terrible password hints, but whatever it was, somebody should lose their Internet		
License!		

### Password guidance

(Modern) NIST guidance	Common guidance
Minimum length >= 8	Min length 8
Maximum length >= 64	Max length 16
Pick and stick*	Change frequently
No algorithmic complexity	Character classes

Some of these pieces of folk wisdom came from the previous NIST password guidance publication, others are accretions of myth on top of legend. The previous version of NIST's guidance was based on a very specific set of assumptions about the environment they'd be used: classified data that needs to be protected for a specific period of time (e.g., 15 years) against a dedicated cracker \_\_\_\_\_\_\_. The new guidance reflects reality for more general computer security uses.

<sup>\*</sup> See Florêncio, Herley and Oorschot, "Password Portfolios and the Finite-Effort User: Sustainably Managing Large Numbers of Accounts", in *USENIX Security 2014*.

# More password guidance

(Modern) NIST guidance	Common guidance
No hints	Allow "rhymes with assword"
No KBA	Non-password passwords
Screen for compromises	???
Careful about 2FA	SMS FTW

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... which will be the topic of our next lecture (2FA)