

# Recall: DAC

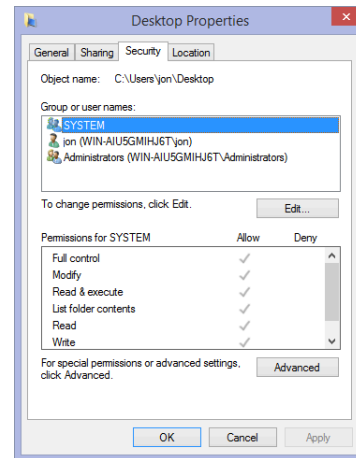
## Organizing principle:

Files and directories have **owners** who have the **discretion** to say who gets to access them.

## Major implementations:

Unix permissions

Access control lists (ACLs)



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This image shows a Windows dialog for editing an access control list (ACL). In this scheme, an owner can write more complex security policies than the cross product of (owner, group, others)  $\times$  (read, write, execute). Instead, the owner can specify a list of arbitrary size, granting or denying permission for any user, group or combination thereof to perform operations on an OS object like a file. ACLs are a very flexible mechanism for representing permissions, but the downside of giving people a lot of flexibility is that they might use it! It's easy to create an ACL so complex that you don't understand it. ACLs are necessary in some circumstances — e.g., centralized file servers that are accessed by thousands of employees from different departments. For a lot of circumstances, however, simple is better.

# MAC: mandatory access control

## Organizing principle:

System administrators can impose access control policies that file owners cannot control or circumvent.

If users can't be trusted... which happens a lot!

# History

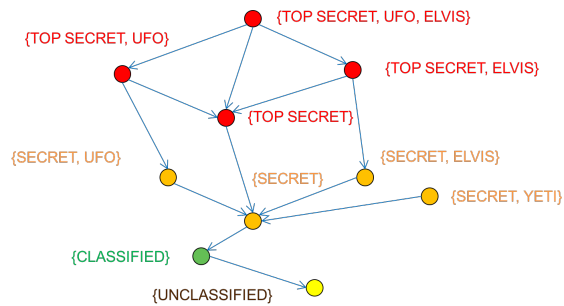
Confidentiality:

"Very hush-hush"

Need-to-know

Formal classification levels\*

Lattices



Source: *Paul Krzyzanowski*

\* See, e.g., the Government of Canada's [Levels of security](#)

† See, e.g., the declassified NSA document [Examples of Lattices](#)

As soon as you start to deal in any quantity of confidential information, it becomes important to describe just \_\_\_\_\_ some information is relative to other information. The phrases "loose lips sink ships" and "it's all very hush-hush" both imply that some information ought to be confidential, but one implies a general disposition towards confidentiality, whereas the other implies that something specific and special requires particular care.

Loose definitions of confidentiality can be described with informal but still well-understood terminology such as "need to know". This approach to maintaining confidentiality is useful, but the human judgement involved is still insufficient for making *automated* decisions about who should be able to access information.

Large, complex organizations that need to specify clear rules for accessing information — traditionally, the military and intelligence services — have created formal classification levels that allow rules to be applied very clearly and definitively, with little judgment required: "is this person cleared to view information with this classification marking?"

Lattices make things even more complex, as they add an \_\_\_\_\_ system of code words.

# The Anderson Report\*†

Trusted computing base

Reference monitor

Policy vs mechanism

So what policies should be enforced?

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\* No relation!

† Anderson, "Computer Security Technology Planning Study", Tech. Rep. ESD-TR-73-51, Vol. II, US Air Force, 1972.

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The Anderson report introduced several key terms and concepts that we rely on today. We've already talked about TCBs, but Anderson also introduced the concept of a *reference monitor*: a system that can \_\_\_\_\_ to information and \_\_\_\_\_ about them. This allows \_\_\_\_\_ to be encoded separately from \_\_\_\_\_: a system provides a "how": how \_\_\_\_\_, and system administrators can supply the "what": \_\_\_\_\_.

# Multi-level security (MLS)

One computer

Many labels

Who can do what to what? *It depends!*

Who can be trusted to specify access control policy?

MAC answer: **System administrators** can impose access control policies that file owners cannot control or circumvent.

# Bell-LaPadula

No read up (confidentiality)

No write down (the *\*-property*)

Administrative burden and high-water marks

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Reference: Bell and LaPadula, "Secure Computer Systems: Mathematical Foundations", *The Mitre Corporation, AD-770 768*, 1973.

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It's a lot of work to label every object in a system. One way to cope with this tsunami of labeling is to allow objects to "float" to the highest label that has written data into them (the "high-water mark"). If a Secret process writes into a Confidential file, instead of disallowing the write, the file can be relabeled as Secret. Thus, any Confidential processes will lose access to the file.

# Biba

Confidentiality not our only goal!

Reads and writes

LOMAC

Windows

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Reference: Biba, "[Integrity Considerations for Secure Computer Systems](#)", *The Mitre Corporation, MTR-3153*, 1975.

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Security isn't just about confidentiality. In some cases, \_\_\_\_\_ of data is more important than its confidentiality. In almost all cases \_\_\_\_\_ is a necessary prerequisite to providing *any* security properties!

We see this used extensively in contemporary operating systems: a process can read from a higher-integrity object (e.g., a file), but not write to them.

LOMAC refers to \_\_\_\_\_ MAC. This is the logical dual of the \_\_\_\_\_ of confidentiality. Mandatory Access Control code that was originally developed for organizations that care about confidentiality can now be used to label objects as "downloaded via a browser", and thus lower-integrity than other files.

Modern version of Windows have four integrity levels: low, medium, high and system. Even if a program is running on behalf of the Administrator, it can't overwrite critical OS files that are labeled with System integrity unless it is itself a System-integrity process (e.g., Windows Update).

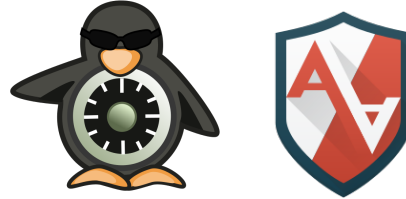
# [Domain and] Type Enforcement

## Categories for subjects and objects

DTE and DTEL

FLASK

SELinux and AppArmor



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Badger, Sterne, Sherman, Walker and Haghghat, "A domain and type enforcement UNIX prototype", *USENIX Comput. Syst.*, vol. 9, no. 1, pp. 47–83, 1996.

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Badger et al.'s *Domain and Type Enforcement* allowed a conventional UNIX machine to be partitioned into various *domains*, and to have MAC policies enforced to constrain the flow of information between them. This included a language for expressing DTE policy (DTEL), and it led to further work on enhancing the security of UNIX and UNIX-like operating systems: [TrustedBSD](#), FLASK, SELinux and AppArmor.



# Linux Security Modules

Patches and problems

"Can you make it a module?"

Comprehensive *hooks* that call arbitrary modules

Separation of *mechanism* from *policy*

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Wright et al., "[Linux Security Modules: General Security Support for the Linux Kernel](#)", in *Proceedings of the 11th USENIX Security Symposium*, 2002.

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This separation of mechanism from policy allows lots of different policies to be enforced, from traditional MAC policies to access control schemes such as Role-Based Access Control and beyond.

# FreeBSD MAC Framework

## Hooks:

```
#ifdef MAC
    error = mac_vnode_check_chdir(td->td_ucred, vp);
    if (error != 0)
        return (error);
#endif
```

## Phones

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Another example of MAC hooks scattered through an operating system is the FreeBSD MAC Framework, which came out of the [TrustedBSD](#) project. Hooks exist to allow a reference monitor to make an access control decision based on a \_\_\_\_\_ (who wants to make the access), an \_\_\_\_\_ (what's being accessed, in this case a file's *vnode* — more about that in ECE 8400 / ENGI 9875) and a currently-installed \_\_\_\_\_ (which may actually be a composition of \_\_\_\_\_).

The FreeBSD MAC Framework is most famously used, not for FreeBSD itself, but to provide a foundation for application sandboxing on iOS and macOS!

# MAC summary

History

MLS

MAC in practice

- Linux Security Modules
- FreeBSD MAC Framework